# Computational Aspects of Mappings

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### (boring!) Definition

Mapping / Map :

A function between domains/spaces

Examples



$$f: \mathbb{R}^2 \to \mathbb{R}^2$$



 $f\colon \mathbb{M}\to \mathbb{R}^2$ 





 $f: \mathbb{M} \to \mathbb{M}'$ 





 $f\colon V\to \mathbb{R}^3$ 



Applications...

### Deformations



### Parameterizations



[Lévy et al. 02]



[Schuler et al. 13]



[Fu et al. 15]





[Weber et al. 12]

### Surface mappings





[Schreiner et al. 04]

[Jin et al. 08]

Discrete maps...

### Smooth case

### e.g., a smooth surface is mapped to $\mathbb{R}^2$



### Discrete case

### e.g., a surface **mesh** is mapped to $\mathbb{R}^2$



# Focus on a Specific embedding

 $\mathbf{x}$ 

 $\Phi_i$ 



1. Boundary mapped to convex polygon





- 1. Boundary constrained to convex polygon
- **2. Discrete Harmonic** Interior vertices at average of neighbors





- 1. Boundary constrained to convex polygon
- 2. Discrete Harmonic Interior vertices at average of neighbors



1. Boundary to convex polygon

2. Harmonic





- 1. Bijective: The graph edges don't overlap themselves
- 2. Discrete Harmonic: analog to smooth harmonic maps



### Continuous harmonic maps

• Harmonic **function**:

 $\Delta f = 0$ 

"Laplacian – difference of value at point to average of neighborhood"



### Continues harmonic maps



- *p* mapped to average of neighborhood!
- By construction, the discrete case















• Imposing constraints

• Finding maps that are most...



### **Constrained Optimization**











 $E(\Phi) = E(A_1, \dots, A_m)$ 

### Map optimization

• In terms of differentials:

 $\operatorname{argmin} E(A_1, \dots, A_m)$ 


# Map optimization

• In terms of differentials:



# Map optimization



#### Must impose continuity!

# Explicit continuity

- Optimization variables:  $A_1, A_2, \dots, A_m$
- Adjacent A<sub>i</sub>'s must agree



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$$A_i v_1 = A_j v_1$$

## Implicit continuity



$$A_{i}[v_{1} v_{2} v_{3}] = [u_{1} u_{2} u_{3}]$$
$$A_{i} = [u_{1} u_{2} u_{3}] [v_{1} v_{2} v_{3}]^{\dagger}$$
$$A_{i} = A_{i}(U)$$

Linearly express  $A_i$ 's in terms of U

## Implicit continuity



• Optimization variables:  $u_1, u_2, \dots, u_n$  (U)

$$E(\Phi) = \sum_{j} f\left(A_{j}(\boldsymbol{U})\right)$$











[Weber & Zorin 2014]



[Weber & Zorin 2014]





[Weber & Zorin 2014]

# Orthogonal and Similarity

• *R* is <u>orthogonal</u> if  $R^T = R^{-1}$ (rotation if det R > 0)

• *S* is a <u>similarity</u> if  $S = \alpha R$ 



# Closest rotation/similarity

- $\mathcal{R}(A)$  = closest orthogonal/rotation matrix to A
- S(A) = closest similarity matrix to A
- Computable using **SVD**:  $A = U\Sigma V^T$ ;  $\Sigma = \text{diag}(\sigma_1, ..., \sigma_n)$

• 
$$\mathcal{R}(A) = U \Sigma V^T = U V^T$$

•  $\mathcal{S}(A) = \overline{\sigma} U V^T$ 

mean of SVs

$$E_L = \sum_j w_j \left\| A_j - \mathcal{S}(A_j) \right\|_F^2$$

#### closest similarity



$$E_L = \sum_j w_j \left\| A_j - \mathcal{S}(A_j) \right\|_F^2$$

#### amount of non-similarity



$$E_L = \sum_j w_j \left\| A_j - S(A_j) \right\|_F^2$$
  
amount of non-similarity

 $E_L = \sum_{j} w_j \left\| A_j - \mathcal{S}(A_j) \right\|_F^2 = 0$ 

#### global similarity = discrete conformal maps



[Lévy et al. 2002]

## As-Rigid-As-Possible (ARAP)

$$E_R = \sum_j w_j \left\| A_j - \mathcal{R}(A_j) \right\|_F^2$$

#### closest rigid transformation



# As-Rigid-As-Possible (ARAP)

$$E_{R} = \sum_{j} w_{j} \left\| A_{j} - \mathcal{R}(A_{j}) \right\|_{F}^{2}$$
amount of non-rigidity

# As-Rigid-As-Possible (ARAP)



[Sorkine & Alexa 2007\*; Chao et al. 2010]

### ARAP vs. LSCM



ARAP vs. LSCM



Dirichlet





LSCM



 $\left\|A_j - \mathcal{S}(A_j)\right\|_{F}^2$ 

ARAP



 $\left\|A_j - \mathcal{R}(A_j)\right\|_F^2$ 

Dirichlet



Least squares

LSCM



 $\left\|A_j - \mathcal{S}(A_j)\right\|_{F}^2$ 

ARAP



 $\left\|A_j - \mathcal{R}(A_j)\right\|_F^2$ 

# Closest Similarity – 2d case

- $\mathcal{S}(A) = \overline{\sigma} U V^T$
- Takes a closed form:

$$A = \begin{bmatrix} a & c \\ b & d \end{bmatrix} = \frac{1}{2} \begin{bmatrix} a+d & c-b \\ b-c & a+d \end{bmatrix} + \frac{1}{2} \begin{bmatrix} a-d & c+b \\ b+c & -a+d \end{bmatrix}$$
  
similarity anti-similarity

Dirichlet



 $\left\|A_j\right\|_F^2$ 

Least squares

LSCM



ARAP



 $\left\|A_j - \mathcal{R}(A_j)\right\|_F^2$ 

Dirichlet



Least squares

LSCM



 $\left\|A_j - \mathcal{S}(A_j)\right\|_{F}^2$ 

2d – Least squares

ARAP



 $\left\|A_j - \mathcal{R}(A_j)\right\|_F^2$ 

Dirichlet



Least squares

LSCM



 $\left\|A_j - \mathcal{S}(A_j)\right\|_{E}^{2}$ 

2d - least squares iterative approximation

ARAP



 $\left\|A_j - \mathcal{R}(A_j)\right\|_F^2$ 

iterative approximation

## Where's the difficulty?

$$E_R = \sum_j w_j \|A_j - \mathcal{R}(A_j)\|_F^2$$

• Not very friendly for direct minimization:

$$A - \mathcal{R}(A) = A - UV^T$$

via SVD of A

But  $\mathcal{R}(A_j)$  is easy to compute...

# Alternating optimization

$$E_R = \sum_j w_j \|A_j - \mathcal{R}(A_j)\|_F^2$$

- Iteratively:

Local step

• Minimize

$$\sum_{j} w_{j} \left\| A_{j} - R_{j} \right\|_{F}^{2}$$
Global step

[Liu et al. 2008]
























## Alternating optimization

#### Very general



#### [Bouaziz et al. 2012]

• Related jargon:

gradient descent, global-local, alternating projections, proximal algorithms

## Singular values perspective

Dirichlet



 $||A||_{F}^{2}$ 

LSCM



 $||A - S(A)||_{F}^{2}$ 

ARAP



 $\|A - \mathcal{R}(A)\|_F^2$ 

## Singular values perspective







 $\sum_k \sigma_k^2$ 



# Singular values



Parameterization





Most Isometric Parameterization [Hormann & Greiner 2000]

Parameterization





[Sorkine et al. 2000]

Parameterization





[Smith & Schaefer 2015]

Surface mapping





[Schreiner et al. 2014]



[Aigerman et al. 2014]

• Volume mapping  $(\sigma_1 - \overline{\sigma})^2 + (\sigma_2 - \overline{\sigma})^2 + (\sigma_3 - \overline{\sigma})^2$ 



[Paillé & Poulin 2012]

• Volumetric mesh improvement





[Freitag & Knupp 2002]

## Spaces of Mappings



## Spaces of Mappings



## Example



As-Rigid-As-Possible

## Spaces of Mappings



## Where's the challenge?

• Singular values = roots of polynomials



#### SV constraints (+energy)



[Lipman 2012]

[Kovalsky et al. 2014]

Approximate via a sequence of convex programs

## Bounding SVs

• Simplest constraint



#### Convex = Simple?

• Cone of positive semidefinite matrices

 $\{A: \mathbf{x}^T A \mathbf{x} \ge 0 \quad \text{for all} \quad \mathbf{x}\}$ 

"easy"

#### Convex = Simple?

Cone of copositive matrices

 $\{A: x^T A x \ge 0 \quad \text{for all} \ x \ge 0\}$ 

"very difficult"

#### Standard convex conic programs

• Linear inequalities



 $\Rightarrow$  linear programming (LP)

#### Standard convex conic programs

• Second order (ice cream) cones

 $\|\mathbf{x}\|_2 \le t$ 



⇒ second order cone programming (SOCP)

#### Standard convex conic programs

• Linear matrix inequalities (LMIs)

 $X \ge 0$ 



 $\Rightarrow$  semidefinite programming (SDP)

## Hierarchy



#### Guarantees & efficient optimization engines!

#### Convexification







## Key observation





**Symmetric** 

Anti-symmetric
### Key observation





**Symmetric** 

Anti-symmetric

### Key observation

$$\gamma \leq \sigma_{\min}(A)$$

$$\gamma \leq \sigma_{\min}\left(\frac{A+A^T}{2}\right)$$

SOCP – 2d SDP – 3d and higher



2d vs. 3d	
2-d	3-d (and higher)
$\mathcal{S}(A)$ has a closed <b>linear</b> form	$\mathcal{S}(A)$ is non-linear

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SOCP	SDP







#### "Most Conformal Mapping"

- Well studied in 2D [Weber et al. 2012]
- Little known in 3D...







# Injectivity

# Injectivity

• "Map is 1-to-1"





### Injective affine map



- Different orientation?
- Not injective on edges



- Injecitivity requires consistent orientation!
  - Is it sufficient?



- Consistent orientation
- Not injective only on inner vertex



• Winding around vertex should be  $2\pi$ 



# Local Injectivity

• "in a small neighborhood, we are injective"



• Note: winding around vertex is always  $2\pi k$ 





[Aigerman2014]

# Local Injectivity

• "in small neighborhood, we are injective"

- Local inj < Global inj</li>
- e.g., f is locally inj

### Important on its own!

# Global injectivity?

an is

- [Tution 1961]: my embedding is injective!
- Soundary doesn't overlap?



# Injective boundary

- Fixed
  - Highly constrained





[Gortler et. al 2006]

# Injective boundary

 Prevent boundary from overlapping during optimization



[Smith and Schaefer 2015]

### Variations on a theme

### Parameterizing a disk...



### Parameterizing a sphere?



### Naïve solution for spheres

Reduce spheres to disks...



### Kind of unnatural





# Instead of fixed boundary...





# Periodic boundary!



### Cone manifolds



[Myles and Zorin 2013]



[Bommes et al. 2009]



[Springborn et al. 2008]



[Aigerman and Lipman 2015]

# Tutte for sphere

#### Linear conditions for bijective parameterization



### 1. Periodic Boundary


### 1. Periodic Boundary

Can glue copies across cuts!



## 1. Periodic Boundary

#### Can (conceptually) tile the whole plane



## 2. Discrete Harmonic Tiling

Each vertex in average of neighbours



## Harmonic tiling of $\mathbb{R}^2$



## Orbifold Tutte Embeddings

lf:

- 1. Boundaries are rotated copies that tile the plane
- 2. The tiling is harmonic **everywhere**

Then:

There exists a unique solution, and it is **injective**!





## Why does it work?

Euclidean orbifold = cone manifold which tiles  $\mathbb{R}^2$ 



## Different cuts yield same embedding



## Embed seamlessly into a "pillow"



#### Surface maps

#### Surface maps

- Input:
  - Two surface meshes M, N
  - Coarse set of corresponding landmarks
- Output: a map  $f: M \to N$ 
  - Bijective (1-1 and onto)
  - High quality (low isometric distortion)
  - Maps landmarks correctly



#### How to represent a surface map?



#### Use flattenings to $\mathbb{R}^2$



#### Cut the mesh and map to disk



#### Flattenings to $\mathbb{R}^2$



#### Recovering the bijection



#### Recovering the bijection



#### Is this good enough?



#### Is this good enough?



#### Is this good enough?



#### Reduce the flattenings' distortion!



#### Let the boundaries move!



#### Let the boundaries move



#### Overlaps





#### Cuts affect mapping!



#### Cuts affect mapping!



#### How to achieve **seamless** result?

![](_page_170_Picture_1.jpeg)

#### Orbifolds are seamless

![](_page_171_Picture_1.jpeg)

#### Define map via orbifold embeddings

![](_page_172_Picture_1.jpeg)

## Summary

# Piecewise linear is simple yet powerful

![](_page_174_Picture_1.jpeg)

## Beyond piecewise linear

• Meshless (e.g., thin plate splines)

![](_page_175_Picture_2.jpeg)

## Beyond piecewise linear

• Higher order FEM

![](_page_176_Picture_2.jpeg)

[Liu et al. 2014]

## What's next?

- Faster optimization
- Coping with non-Euclidean domains
- New distortion metrics

## THE END

![](_page_178_Picture_1.jpeg)

#### (Some things cannot be mapped)